Probability Bounds on Persistent Centroid Appearance in Random Growing Trees

Ryan Hu

St. John's College

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• Network archaeology (Haigh, 1970; Shah & Zaman, 2011; Bubeck, Devroye & Lugosi, 2017)

Definition (Uniform Attachment Model)

A family of random trees (i.e. family of rv's whose values are trees) $\{T_n\}_{n\geq 1}$ is said to follow a uniform attachment model if $|T_n| = n$, $V(T_n) \subseteq V(T_{n+1})$ for all $n \in \mathbb{N}$, and, if we write $V(T_n) = \{v_1, \ldots, v_n\}$, then $\mathbb{P}((v_i, v_{n+1}) \in E(T_{n+1})) = 1/n$ for each $i \leq n$. Also, each time the attachment choices should be independent from the previous ones.

Definition (Centrality)

Given a finite tree T and a vertex $v \in V(T)$. Consider the forest formed by deleting this vertex. The centrality of this vertex is then

 $\psi_T(\mathbf{v}) :=$ The Size of the Largest Tree in the Forest Formed.

Definition (Centroid)

Given a finite tree T and a $v \in V(T)$. v is called a centroid of T if it minimises $\psi_T(v)$.

Theorem (Loh, Jog, 2016)

Let (T_n) be a family of random growing trees given by the uniform attachment. Then, with probability 1, there exists a $N \in \mathbb{N}$ and a $v \in \bigcup_n V(T_n)$ such that for all $n \ge N$, v is the centroid of T_n .

Definition

Let (T_n) be a family of random growing trees given by uniform attachment. We define the random variable

 $T := \min\{N \in \mathbb{N} : \text{there exists } k \in \mathbb{N} \text{ s.t. } C(T_n) = \{v_k\} \text{ for all } n \ge N\}.$

• Note that by the above theorem, this random variable is almost surely well-defined.

Probability Bounds on Persistent Centroid Appearance

Goal

Our goal is to find an explicit expression $N(\varepsilon)$ such that

 $\mathbb{P}(T > N(\varepsilon)) \leq \varepsilon.$

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Lemma

Let T be a finite tree. Then T has at most 2 centroids. And, if T has 2 centroids, then these 2 centroids are connected by an edge.

• Start by noting that, whenever there is a change in centroid, there must be a stage in between at which there's two centroids. Hence,

$$\{T > N\} = \bigcup_{k=N}^{\infty} \{|\mathcal{C}(T_k)| = 2\}.$$

• The union bound of the above is not enough, instead, consider the union bound given by

$$\{T > N\} = \bigcup_{j=2}^{\infty} \{\text{there exists } k \ge N, i < j \text{ s.t. } \mathcal{C}(T_k) = \{v_i, v_j\}\}.$$

• It's easy to track $(\psi_{T_k}(v_i), \psi_{T_k}(v_j))$ as k evolves, it's just a Polya urn!

Definition

A sequence of random variables (X_i) is said to be exchangeable if $(X_1, \ldots, X_n) \stackrel{d}{=} (X_{\pi(1)} \ldots, x_{\pi(n)})$ for all $n \in \mathbb{N}$.

Theorem (de Finetti, 1931)

A sequence of Bernoulli random variables $(X_i)_{i\geq 1}$ is exchangeable if and only if there exists a distribution function F on [0,1] such that for all $n \in \mathbb{N}$,

$$\mathbb{P}(X_1 = x_1, \ldots, X_n = x_n) = \int_0^1 \theta^{\sum_{i=1}^n x_i} (1-\theta)^{n-\sum_{i=1}^n x_i} \, dF(\theta).$$

Definition (Polya Urn)

A Polya urn puts x white and y black balls into an urn. At each step, one ball is drawn uniformly at random from the urn, and its color observed; it is then returned in the urn, and an additional ball of the same color is added to the urn.

Theorem (Janson, 2019)

Let (X_n, Y_n) be the Polya urn model with $X_0 = a, Y_0 = b$. Then, there exists a rv $W \sim \text{beta}(a, b)$ with $\frac{X_n}{X_n + Y_n} \to W$ almost surely. And, for each $w \in [0, 1]$, conditioning on W = w, $X_n = X_0 + \sum_{k=1}^n Q_k$ with $Q_k \stackrel{\text{i.i.d.}}{\sim}$ Bernoulli(w).

Law of Iterated Logarithm Bounds

 Law of Iterated Logarithm is a classical result in probability theory. It says

$$\limsup_{n \to \infty} \frac{\left|\sum_{k=1}^{n} Y_{k}\right|}{\sqrt{2n \log \log n}} = 1 \text{ a.s.},$$

where (Y_k) are i.i.d. rvs with zero mean and unit variances.

• Recent results in sequential hypothesis testing use a non-asymptotic result (Balasubramani & Ramdas, 2016):

Theorem

Let $Q_k \overset{\text{i.i.d.}}{\sim} \text{Bernoulli}(w)$, $\eta_t = \eta_0 + \sum_{k=1}^t (-1)^{Q_k}$. Then, with probability $\geq 1 - \delta$, for all $t \geq \frac{C_1}{w(1-w)} \log(4/\delta)$ uniformly, we have

 $|\eta_t - (2w-1)t| \leq \sqrt{C_2w(1-w)t(2\log\log(C_3w(1-w)t/|\eta_t|) + \log(2/\delta))},$

where C_1, C_2, C_3 are some explicit constants.

- Analyzed time to persistence for centroid in uniform attachment trees.
- Results rely on convergence in Polya urns (de Finetti's Theorem), finite-time Law of the Iterated Logarithm.
- Some details still to be worked out.
- Future work: Preferential attachment trees, diffusion in d-regular trees.